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A Fourth-Order Nonlinear Conjugate Gradient Method in Equality Constrained Optimization

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Abstract

This paper presents a fourth-order nonlinear conjugate gradient method in equality constrained optimization. The idea is to transform the constrained problem into unconstrained type through the Lagrange multipliers scheme. Using four terms of Taylor series development, we approximate the transformed function (augmented Lagrange function). Lastly, we employ the new fourth-order nonlinear conjugate gradient method in equality constrained optimization to solve the optimization problem. We present the algorithm in steps and some properties of the gradients are proved, using classical results. Also, the convergence analysis has been proved under classical and known assumptions. Furthermore, we present the obtained numerical results and compare them to some existing results. The analysis of results confirms that the new method is accurate.

Keywords: Fourth-order conjugate gradient method; equality constrained optimization; objective function; nonlinear polynomial approximation; Lagrange multipliers scheme.

Mathematical subject classification (2010): 65K10.

1 Introduction

The equality constrained optimization of a smooth function, Q, in many variables remains an important problem in optimization theory. This is true since many scientists seek to solve this class of problems, in real life applications. Every equality constrained optimization problem [1] could be put in the form

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$$Optimize_{x} Q(x) \tag{1}$$

subject to

$$w_i(x) = b_i, i = 1, 2, ..., m,$$

where $Q(x) \in \Re$, $W_i(x) \in \Re^m$ and $b_i \in \Re$. The general approach is to transform the problem into unconstrained type by Lagrange multipliers scheme and solve the zeros of the function gradient since the local minima occur at stationary points. A fourth-order nonlinear conjugate gradient method in equality constrained optimization finds the global minimum of the transformed function. The new method is characterized by the following. Consider the transformed case of problem (1):

$$\min_{x \in \Re^n} f(x) \tag{2}$$

where f is a differentiable function. We note that $\max_{x \in \Re^n} f(x) = \min_{x \in \Re^n} (-f(x))$. In order to solve this unconstrained problem, we need to design a special algorithm that reduces the high storage and computation cost of some computed matrices [2]. Various types of conjugate gradient method have been used to solve unconstrained minimization problems [3]. Usually, a function F is constructed to approximate f. If the objective function is not quadratic or the inexact line search is used, some of the conjugate gradient methods fail to converge globally [4,5]. The process of minimizing a non-quadratic objective function through the conjugate gradient method is called the nonlinear conjugate gradient method [6,7]. Many scholars have published their findings on this method [8,9,10]. New algorithms on nonlinear conjugate gradient method are available [11,12,13,14,15]. Every conjugate gradient method is an iterative scheme of the form

$$x_{k+1} = x_k + \alpha_k d_k, \quad k = 0, 1, 2, \dots$$
 (3)

where x_0 is an initial point, α_k is a step size and the search direction

$$d_{k} = \begin{cases} -g_{k}, k = 0, \\ -g_{k} + \beta_{k}d_{k}, k > 0. \end{cases}$$
(4)

 $g_k = \nabla f(x_k)$ and β_k specifies the choice of conjugate gradient method [15]. It could take any of the following forms.

$$\boldsymbol{\beta}^{FR} = \frac{\|\boldsymbol{g}_{k}\|^{2}}{\|\boldsymbol{g}_{k-1}\|^{2}} [5], \ \boldsymbol{\beta}^{PRP} = \frac{\boldsymbol{g}_{k}^{T}(\boldsymbol{g}_{k}-\boldsymbol{g}_{k-1})}{\|\boldsymbol{g}_{k-1}\|^{2}} [17], \ \boldsymbol{\beta}^{HS} = \frac{\boldsymbol{g}_{k}^{T}(\boldsymbol{g}_{k}-\boldsymbol{g}_{k-1})}{\boldsymbol{d}_{k-1}^{T}(\boldsymbol{g}_{k}-\boldsymbol{g}_{k-1})} [18,19], \\ \boldsymbol{\beta}^{CD} = -\frac{\|\boldsymbol{g}_{k}\|^{2}}{\boldsymbol{d}_{k-1}^{T}\boldsymbol{g}_{k-1}} [15], \ \boldsymbol{\beta}^{DY} = \frac{\|\boldsymbol{g}_{k}\|^{2}}{\boldsymbol{d}_{k-1}^{T}(\boldsymbol{g}_{k}-\boldsymbol{g}_{k-1})} [19], \ \boldsymbol{\beta}^{LS} = -\frac{\boldsymbol{g}_{k}^{T}(\boldsymbol{g}_{k}-\boldsymbol{g}_{k-1})}{\boldsymbol{d}_{k-1}^{T}\boldsymbol{g}_{k-1}} [20]$$

and

Many of these conjugate gradient methods use inexact line search technique [21]. Others use exact line search approach [22]. Stoer [23] studied the conjugate gradient method on a subspace and obtained a variant of the method with an inexact line search approach. The search for a reliable and accurate scheme motivated this work on a fourth-order nonlinear conjugate gradient method (FONCGM) in equality constrained

optimization. This method is presented in seven sections. Sections (two and three) discuss the transcription of the equality constrained problem to unconstrained type and FONCGM, respectively. In section four, we give the convergence analysis. Section five presents some test problems. Section six explains the numerical results while section seven ends this work with a conclusion.

2 Transcription of Equality Constrained Optimization Problem

We transform the constrained problem (1) into unconstrained problem of the form

Optimize
$$f(x, \lambda)$$
 (5)

 $x \in \Re^N$, $\lambda \in \Re^m$ and f is the Lagrange's function defined by

 $f(x,\lambda) = Q(x) + \sum_{i=1}^{m} \lambda_i (b_i - W_i(x))$. λ_i , i = 1, 2... m are the Lagrange multipliers. A solution to

problem (1) can then be found by solving problem (5) if there exists a vector

 $\lambda = (\lambda_{1}, \lambda_{2}, ..., \lambda_{m})$ such that

$$\frac{\partial f(x,\lambda)}{\partial x_1} = 0, \ \frac{\partial f(x,\lambda)}{\partial x_2} = 0, \ \dots, \ \frac{\partial f(x,\lambda)}{\partial x_N} = 0 \tag{6}$$

and

$$\frac{\partial f(x,\lambda)}{\partial \lambda_1} = 0, \quad \frac{\partial f(x,\lambda)}{\partial \lambda_2} = 0, \quad \dots, \quad \frac{\partial f(x,\lambda)}{\partial \lambda_m} = 0. \tag{7}$$

Equations (6) and (7) will generate a set of N + m equations in N + m unknowns to be solved. As the dimensionality of the problem increases, we have many equations to solve simultaneously. The problem computation becomes very tedious, if analytic method is employed. Thus, we develop and apply the fourth-order nonlinear conjugate gradient method where the objective function is the augmented Lagrange function defined by

$$f(x,\lambda,\mu) = Q(x) + \sum_{i=1}^{m} \lambda_i (b_i - W_i(x)) + \frac{1}{2} \sum_{i=1}^{m} \mu (b_i - W_i(x))^2$$
(8)

 μ is a scalar called the penalty parameter. The multipliers updates usually take the form

$$\lambda_{i+1} = \lambda_i + \mu(b_i - W_i(x)) \tag{9}$$

Our approach is to choose initial vectors x_0, λ_0 , a parameter μ_k and use the new fourth-order nonlinear conjugate gradient method to optimize $f(x, \lambda, \mu_k)$ over \Re^N . The scalars $\mu_0, \mu_1, ..., \mu_k$ could be determined on the basis of results obtained during iteration process. The new fourth-order nonlinear conjugate gradient method in equality constrained optimization follows.

3 The Fourth-Order Nonlinear Conjugate Gradient Method (FONCGM)

The fourth-order nonlinear conjugate gradient method is based on four terms Taylor series representation of f by F. This representation is expected to be a better approximation of f than the usual representation. The following is the representation of F at point x_k .

$$F(x) = f(x_k) + df(x_k) + \frac{1}{2!}d^2 f(x_k) + \frac{1}{3!}d^3 f(x_k) + \frac{1}{4!}d^4 f(x_k),$$
(10)

where

$$d^{n}f(x_{k}) = \sum_{i_{1}}^{N} \sum_{i_{2}}^{N} \dots \sum_{i_{n}}^{N} h_{i_{1}}h_{i_{2}}\dots h_{i_{n}} \frac{\partial^{n}f(x_{k})}{\partial x_{i_{1}}\partial x_{i_{2}}\dots \partial x_{i_{N}}}, x, x_{k}, h_{i_{j}} \in \Re^{N}; h_{i_{j}} = x - x_{k}, \ 2 \le n \le 4.$$

Using a vector $h = x - x_k$ and $A_i = \nabla^i f(x_k)$, in equation (10), we have

$$F(x) = f(x_k) + h^T A_1 + \frac{1}{2!} h^T A_2 h + \frac{1}{3!} h^T (h^T A_3 h) + \frac{1}{4!} h^T (h^T A_4 h) h$$

= $f(x_k) + h^T A_1 + \frac{1}{2!} h^T \left[A_2 + \left\{ \frac{2}{3!} A_3 + \frac{2}{4!} A_4 h \right\} h \right] h$ (11)

Using tensor notations presented in [24], we have

$$\begin{bmatrix} \frac{2}{3!}A_3 + \frac{2}{4!}A_4h \end{bmatrix} h = \begin{bmatrix} \sum_{j=3}^4 \frac{2}{j!}A_j \prod_{p=1}^{j=3} h^{Z_p} \\ = \begin{bmatrix} \sum_{j=3}^4 \frac{2}{j!} \sum_{m=0}^{j=3} (-1)^m \binom{j-3}{m} g(x_k + \{(j-3) - m\}h) \end{bmatrix}^T h$$
$$= \frac{1}{12} [g(x) + 3g(x_k)]^T h$$
(12)

where g(x) denotes the gradient of f, at point x, $\prod_{p=1}^{2} h^{Z_p} = h^T h$, $\binom{n}{m} = \frac{n!}{(n-m)!m!}$, $Z_p = T^{\frac{1}{2}\left[1 + (-1)^p\right]}$ and T denotes the gradient of f.

T denotes transpose. It follows that

$$F(x) = f(x_k) + h^T A_1 + \frac{1}{2} h^T H(x) h$$
(13)

where

$$H(x) = A_2 + \frac{1}{12} [g(x) + 3g(x_k)]^T h.$$

Similarly,

$$\nabla F(x) = A_1 + A_2 h^T + \frac{1}{2} h^T A_3 h + \frac{1}{3!} h^T (h^T A_4 h)$$

$$= A_1 + \sum_{j=2}^{4} \frac{1}{(j-1)!} A_j \prod_{p=1}^{j-1} h^{Z_p}$$

$$= A_1 + \sum_{j=2}^{4} \frac{1}{(j-1)!} \sum_{m=0}^{j-1} (-1)^m {j-1 \choose m} g[x_k + \{(j-1)-m\}h]$$

$$= A_1 + \frac{1}{6} [g(x_k + 3h) + 3g(x_k + h) - 4g(x_k)]$$

$$= A_1 + \frac{1}{6} [g(3x - 2x_k) + 3g(x) - 4g(x_k)]$$
(14)

$$\nabla F(x_{k+1}) = A_1 + \frac{1}{6} \left[g(3x_{k+1} - 2x_k) + 3g(x_{k+1}) - 4g(x_k) \right]$$
(15)

$$\nabla F(x_k) = A_1 = g(x_k). \tag{16}$$

$$\nabla F(x_{k+1}) = \nabla F(x_k) + \frac{1}{6} \Big[g(3x_{k+1} - 2x_k) + 3g(x_{k+1}) - 4g(x_k) \Big]$$

= $\frac{1}{6} \Big[g(3x_{k+1} - 2x_k) + 3g(x_{k+1}) - 4g(x_k) + 6g(x_k) \Big]$
= $\frac{1}{6} \Big[g(3x_{k+1} - 2x_k) + 3g(x_{k+1}) + 2g(x_k) \Big].$ (17)

Using $G_{k+1} = \nabla F(x_{k+1})$, we present a fourth-order nonlinear conjugate gradient algorithm in which the directions of search, $D_0, D_1, ..., D_k$ are H conjugate. That is,

$$D_{k+1}^T H_{k+1} D_k = 0 (18)$$

From the classical results, it follows that

$$G_{k+1} = G_k + \alpha_k H_{k+1} D_k \tag{19}$$

With a given x_0 , α_k is computed such that

$$F(x_{k} + \alpha_{k}D_{k}) < F(x_{k}),$$

$$D_{k+1} = \begin{cases} -G_{k+1}, & k = 0 \\ -G_{k+1} + \beta_{k}D_{k}, & k > 0 \end{cases}$$
(20)

and

$$x_{k+1} = x_k + \alpha_k D_k$$
, $k = 0, 1, ...$

From equations (12) and (15),

$$\begin{bmatrix} -G_{k+1} + \beta_k D_k \end{bmatrix}^T H_{k+1} D_k = 0; \ \beta_k D_k^T H_{k+1} D_k = G_{k+1}^T H_{k+1} D_k. \text{ From equation (19),} \\ \beta_k = \frac{G_{k+1}^T (G_{k+1} - G_k)}{D_k^T (G_{k+1} - G_k)}$$
(21)

The algorithm is described below.

Algorithm 1. (FONCGM)

Step 1: Select $x_0 \in \mathfrak{R}^N$, $N \ge 2$, $\|.\|$ is Euclidean norm, λ_0 , μ_k and $\mathcal{E} > 0$ (a small number: 0.000001). Set $G_0 = \nabla f_{\mu}(x_0, \lambda_0, \mu_k)$, $D_0 = -G_0$ and k = 0. Step 2: If $\|G_k\| \le \varepsilon$, stop. Choose x_k , otherwise go to step 3. Step 3: Compute α_k such that $F(x_k + \alpha_k D_k) < F(x_k)$ and go to step 4. Step 4: Compute

$$G_{k+1} = \frac{1}{6} \left[g(3x_{k+1} - 2x_k) + 3g(x_{k+1}) + 2g(x_k) \right] \beta_k = \frac{G_{k+1}^T (G_{k+1} - G_k)}{G_k^T (G_{k+1} - G_k)}, D_{k+1} = -G_k + \beta_k D_k.$$

$$x_{k+1} = x_k + \alpha_k D_k.$$

Go to step 5.

Step 5: Check for optimality of G.

 $\text{If } G_{i+1} = 0 \ and \ \mid \lambda_{i+1} - \lambda_i \mid > \varepsilon \ \text{, go to step 6.If } G_{i+1} = 0 \ and \ \mid \lambda_{i+1} - \lambda_i \mid \le \varepsilon \ \text{, go to step 7.}$

Step 6: Update λ as follows.

$$\lambda_{i+1} = \lambda_i + \mu_k (b_i - W_i(x_i))$$

Choose $\mu_{k+1} > \mu_k$ such as $\mu_{k+1} = \mu_k + 0.1 * 0.2^j$; $j \le 1000$; $\mu_0 = 0$

Set k = k + 1. Go to step 2.

Step 7: Stop iteration. x_{i+1} is the final optimal point found

Remark: Dai and Yuan [12] presented a nonlinear conjugate gradient algorithm for solving unconstrained optimization problems. Below is Dai-Yuan's algorithm for problem (1).

Algorithm 2. (Nonlinear conjugate gradient method)

Step 1: Select $x_0 \in \Re^N$, $N \ge 2$ and $\mathcal{E} > 0$. Set $d_0 = -g_0$ and k = 0. Step 2: If $||g_k|| \le \varepsilon$, stop. Take x_k . Otherwise go to step 3. Step 3: Compute α_k such that $f(x_k + \alpha_k d_k) < f(x_k)$, go to step 4. Step 4: Compute $\beta_k = \frac{||g_k||^2}{d_{k-1}^{T}(g_k - g_{k-1})}$, $d_{k+1} = -g_k + \beta_k d_k$ and $x_{k+1} = x_k + \alpha_k d_k$. Step 5: Set k = k+1. Contacting 2.

Step 5: Set k = k+1. Go to step 2.

4 Convergence Analysis

We employ the convergence results of algorithm (2), as contained in the following lemma and theorem, to establish the convergence of algorithm (1). We assume that the objective function satisfies the following conditions.

4.1 Assumptions

- I. f is bounded below in \mathfrak{R}^N and is four times continuously differentiable in a neighborhood Zof the level set $L = \{x \in \mathfrak{R}^N : f(x) \le f(x_0)\}$
- II. The gradient, g(x), is Lipschitz continuous in Z, namely, there exists a constant Lc > 0 such that $\|\nabla f(x) \nabla f(y)\| \le Lc \|x y\|$, $x, y \in Z$.
- III. The extended hessian matrix H(x) is positive definite.

4.2 Lemma

I. Suppose that x_0 is a starting point for which the above assumptions are satisfied. Consider any method of the form (2), where D_k , a vector, is the descent direction and α_k satisfies the standard Wolfe conditions [18], then

$$\sum_{k\geq 0} \frac{\left(G_k^T D_k\right)^2}{\parallel D_k \parallel^2} < \infty$$

II. Suppose that x_0 is a starting point for which the above assumptions are satisfied. Let $\{x_k, k = 1, 2, ...\}$ be generated by algorithm (1). Then, the algorithm either terminates at a stationary point or converges in the sense that

$$\lim_{k\to\infty}\inf \parallel G(x_k)\parallel=0$$

III.

Theorem 1. Suppose that f is continuously differentiable, bounded below and the norm of the Hessian matrix is bounded. The iteration $\{x_k\}$ is generated by algorithm (2) satisfies $x_k \to x^*$ as $k \to \infty$ and the Hessian matrix of f is positive definite. Let ε_k be the relative error in the truncated conjugate gradient method and the algorithm. If $\varepsilon_k \to \infty$ then $\{x_k\}$ converges, that is,

$$\lim_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0$$

Proof (Lemma (i)):

Dai and Yuan proved this lemma for algorithm (2): $\sum_{k\geq 0} \frac{(g_k^T d_k)^2}{\|d_k\|^2} < \infty$. Similar proof is presented for algorithm (1), since $G_k = G(x_k) = \nabla F(x_k) = g(x_k)$) and the search directions d_k and D_k have same definition. This

is obvious on using the assumptions of lemma (i), Dai and Yuan's proof and k in place of k + 1 in equation (17).

$$\begin{split} & \frac{\left(G_{k}^{T}D_{k}\right)^{2}}{\|D_{k}\|^{2}} = \frac{\frac{1}{36}\left\{g\left(3x_{k}-2x_{k-1}\right)+3g\left(x_{k}\right)+2g\left(x_{k-1}\right)\right\}^{T}D_{k}\right\}^{2}}{\|D_{k}\|^{2}} \\ &= \frac{\frac{1}{36}\left\{g\left(3x_{k}-2x_{k-1}\right)+3g\left(x_{k}\right)+2g\left(x_{k-1}\right)\right\}^{T}d_{k}\right\}^{2}}{\|d_{k}\|^{2}}, d_{k} = D_{k} \\ &= \frac{\frac{1}{36}\left\{g\left(3x_{k}-2x_{k-1}\right)^{T}d_{k}\right)+3\left(g\left(x_{k}\right)^{T}d_{k}\right)\right\}^{2}}{\|d_{k}\|^{2}}, g\left(x_{k-1}\right)^{T}d_{k} = 0 \text{ (by choice of } \alpha_{k}) \\ &\leq \frac{\frac{1}{36}\left\{g\left(3x_{k}-2x_{k-1}\right)^{T}d_{k}\right)^{2}+9\left(g\left(x_{k}\right)^{T}d_{k}\right)^{2}+6\left(g\left(3x_{k}-2x_{k-1}\right)^{T}d_{k}\right)^{2}+6\left(g\left(x_{k}\right)^{T}d_{k}\right)^{2}\right\}}{\|d_{k}\|^{2}} \\ &= \frac{1}{36}\left\{\frac{\left(g\left(3x_{k}-2x_{k-1}\right)^{T}d_{k}\right)^{2}}{\|d_{k}\|^{2}}+9\frac{\left(g\left(x_{k}\right)^{T}d_{k}\right)^{2}}{\|d_{k}\|^{2}}+6\frac{\left(g\left(3x_{k}-2x_{k-1}\right)^{T}d_{k}\right)^{2}}{\|d_{k}\|^{2}}\right\}}{\sum_{k\geq 0}\frac{\left(G_{k}^{T}D_{k}\right)^{2}}{\|D_{k}\|^{2}}\leq \frac{1}{36}\left\{\sum_{k\geq 0}\frac{\left(g\left(3x_{k}-2x_{k-1}\right)^{T}d_{k}\right)^{2}}{\|d_{k}\|^{2}}+9\sum_{k\geq 0}\frac{\left(g\left(x_{k}\right)^{T}d_{k}\right)^{2}}{\|d_{k}\|^{2}}+6\sum_{k\geq 0}\frac{\left(g\left(x_{k}\right)^{T}d_{k}\right)^{2}}{\|d_{k}\|^{2}}\right\}<\infty \end{split}$$

Thus,

$$\sum_{k\geq 0} \frac{\left(G_k^T D_k\right)^2}{\|D_k\|^2} < \infty \text{ as required}$$

Proof (Lemma (ii)):

Dai and Yuan proved this lemma for algorithm (2). The proof is same since, from equation (16), $G_k = G(x_k) = \nabla F(x_k) = g(x_k)$. It follows that

 $\lim_{k \to \infty} \inf \| G_k \| = \lim_{k \to \infty} \| g_k \| = 0.$ This is true since

$$\lim_{k \to \infty} \inf \|G_{k+1}\| = \lim_{k \to \infty} \inf \frac{1}{6} \| \{ g(3x_{k+1} - 2x_k) + 3g(x_{k+1}) + 2g(x_k) \} \|$$

$$\leq \liminf_{k \to \infty} \inf \frac{1}{6} \| g(3x_{k+1} - 2x_k) \| + \liminf_{k \to \infty} \inf \frac{1}{6} \| 3g(x_{k+1}) \| + \liminf_{k \to \infty} \inf \frac{1}{6} \| 2g(x_k) \| = 0$$

 $\lim_{k\to\infty} \inf \|G_{k+1}\| \le 0 \text{ . But } \|G_{k+1}\| \ge 0 \text{ implies that } \lim_{k\to\infty} \inf \|G_{k+1}\| = 0 \text{ or } \lim_{k\to\infty} \inf \|G_k\| = 0 \text{ as required.}$

Proof of theorem (1): The proof is available in many literatures. Noting that $G_k = G(x_k) = \nabla F(x_k) = g(x_k)$, the proof is same since the assumptions on algorithm (1) meet the requirements of this theorem. Using $M \ge m \ge 0$, $\varepsilon_k = x_k - x^*$ and the results from NMC [25], we have

$$\begin{split} \| x_{k+1} - x^* \|^2 &= (x_{k+1} - x^*)^T (x_{k+1} - x^*) \\ &= (x_{k+1} - x^*)^T (x_k + \alpha_k d_k - x^*) \\ &\leq \frac{M}{n} \left(\frac{1 - R}{1 + R}\right)^{2(k+1)} \| x_k - x^* \|^2 \leq \frac{M}{n} \left(\frac{1 - R}{1 + R}\right)^{2(k+1)} \| x_0 - x^* \|^2; \ R &= \frac{m}{M}. \\ &\| x_{k+1} - x^* \| \leq \sqrt{\frac{M}{n}} \left(\frac{1 - R}{1 + R}\right)^{(k+1)} \| x_k - x^* \| \\ &\lim_{k \to \infty} \frac{\| x_{k+1} - x^* \|}{\| x_k - x^* \|} = 0, \text{ since } \lim_{k \to \infty} \sqrt{\frac{M}{n}} \left(\frac{1 - R}{1 + R}\right)^{(k+1)} = 0. \end{split}$$

End of proof.

5 Numerical Consideration

To illustrate the behavior of the algorithm proposed in this paper, we wrote MATLAB codes for solving the following problems and ran them on a PC with Windows 7. The gradient tolerance is 0.000001. The problems are of the form (1) with the following expressions for Q(x) and the constraints.

Problem 1 (HS48 [26])

$$Q(x) = (x_1 - 1)^2 + (x_2 - x_3)^2 + (x_4 - x_5)^2$$

such that $x_1 + x_2 + x_3 + x_4 + x_5 - 5 = 0$

$$x_{3} + 2(x_{4} + x_{5}) + 3 = 0$$

$$x_{0} = (3, 5, -3, 2, -2); \quad x^{*} = (1, 11/3, 11/3, -5/3, -5/3); \quad f(x^{*}) = 0$$

Problem 2 (HS51 [26])

$$Q(x) = (x_1 - x_2)^2 + (x_2 + x_3 - 2)^2 + (x_4 - 1)^2 + (x_5 - 1)^2$$

such that $x_1 + 3x_2 - 4 = 0$

$$x_3 + x_4 - 2x_5 = 0$$

 $x_2 - x_5 = 0$
 $x_0 = (2.5, 0.5, 2, -1, 0.5); \ x^* = (1, 1, 1, 1, 1); \ f(x^*) = 0$

Problem 3 (HS50 [26])

$$Q(x) = (x_1 - x_2)^2 + (x_3 - 1)^2 + (x_4 - 1)^4 + (x_5 - 1)^6$$

such that $x_1 + x_2 + x_3 + 4x_4 - 7 = 0$

$$x_3 + 5x_5 - 6 = 0$$

 $x_0 = (10, 7, 2, -3, 0.8); \ x^* = (1, 1, 1, 1, 1); \ f(x^*) = 0.$

Problem 4 [27], problem (11)

$$Q(x) = \sum_{i=1}^{100} i x_i^4$$

such that $c_i(x) = \sum_{j=1}^{i+1} x_j - \frac{i}{10} = 0$, i = 1, 2, ..., 20. $x_0 = (0.25, 0.25, ..., 0.25); x^* = (0.0557, 0.0443, ..., -0.0011, -0.0011); f(x^*) = 0.0228$.

Problem 5 [28], problem (394)

$$Q(x) = \sum_{i=1}^{20} i(x_i^2 + x_i^4)$$

such that $c_1(x) = \sum_{i=1}^{20} x_i^2 - 1 = 0$ $x_0 = (2, 2, 2, ..., 2); x^* = (0.91287, 0.408268, -0.000017, ..., -0.0000014); f(x^*) = 1.91667.$

Problem 6 [29], problem 8

$$Q(x) = \frac{1}{2} \sum_{i=1}^{k-2} (x_{k+i+1} - x_{k+i})^4$$

such that $x_{k+i} - x_{i+1} + x_i = i$, i = 1, 2, ..., k - 1.

$$x_0 = (1, 2, ..., k, 2, 3, ..., k); k = 3; f(x^*) = 0.$$

Problem 7 [29], problem 8

$$Q(x) = \frac{1}{2} \sum_{i=1}^{k-2} (x_{k+i+1} - x_{k+i})^4$$

such that $x_{k+i} - x_{i+1} + x_i = i$, i = 1, 2, ..., k - 1.

$$x_0 = (1, 2, ..., k, 2, 3, ..., k); k = 10; f(x^*) = 0.$$

Problem 8 [29], problem 10

$$Q(x) = \sum_{i=1}^{n} \cos^2 (2\pi x_i \sin(\pi / 20))$$

such that $x_i - x_{i+1} = 0.4$, i = 1, 2, ..., n-1.

$$x_0 = (1, 0.6, 0.2, ..., 1 - 0.4((i - 1)); n = 3; f(x^*) = 0.$$

Problem 9 [29], problem 10

$$Q(x) = \sum_{i=1}^{n} \cos^2 (2\pi x_i \sin(\pi / 20))$$

such that $x_i - x_{i+1} = 0.4$, i = 1, 2, ..., n-1.

$$x_0 = (1, 0.6, 0.2, ..., 1 - 0.4((i-1)); n = 10; f(x^*) = 0.$$

Problem 10 [29], problem 9

$$Q(x) = \sum_{i=1}^{k-2} 100(x_{k+i+1} - x_{k+i})^2 + (1 - x_{k+i})^2$$

such that $x_{k+i} - x_{i+1} + x_i = i$, i = 1, 2, ..., k - 1.

$$x_0 = (1, 2, ..., k, 2, 3, ..., k); k = 3; f(x^*) = 0.$$

Problem 11 [29], problem 9

$$Q(x) = \sum_{i=1}^{k-2} 100(x_{k+i+1} - x_{k+i})^2 + (1 - x_{k+i})^2$$

such that $x_{k+i} - x_{i+1} + x_i = i$, i = 1, 2, ..., k - 1.

$$x_0 = (1, 2, ..., k, 2, 3, ..., k); k = 10; f(x^*) = 0.$$

Problem 12 [30], Problem 4

$$Q(x) = e^{x_1 x_2 x_3 x_4 x_5}$$

such that $x_1^2 + x_2^2 + x_3^2 + x_4^2 + x_5^2 - 10 = 0$

$$x_{2}x_{3} - 5x_{4}x_{5} = 0$$

$$x_{1}^{3} + x_{2}^{3} + 1 = 0$$

$$x_{0} = (-2, 2, 2, -1, -1); \quad f(x^{*}) = 0.05395$$

The numerical results obtained for the new method vis-a-vis some classical methods (FOCGM, FR, DY, PRP) are presented in Table 1, (P: Problem; ITE: Number of iterations; TIME: Computer execution time (s); FN.: Function value at the end of iterations). The stopping criterion is $||g(x_k)|| < 0.000001$ while the maximum number of iterations is 1000.

Performance profiles have been introduced by Dolan and More' [31]. The main idea is to show, graphically, the performance of various solvers on a given set of problems. That is, the curves are used to compare the efficiency of a set S of solvers on a set P of test problems. $t_{p,s}$ denotes the performance of a solver s (based on the number of iterations, function evaluations, gradient evaluations or execution time) on the problem p. $r_{p,s}$ denotes the relative performance of a solver s on a problem p and

 $r_{p,s} = \frac{t_{p,s}}{\min\{t_{p,s} : s \in S\}}.$ Our assumption is that $r_{p,s} \le w, w \in \Re$, for all solvers s on the problems p.

$$r_{p,s} = w$$
 if solver s cannot solve problem p. The performance profile of the solver s is the function

$$y_s : [1, w] \to [0, 1]$$
 such that $y_s(t) = \frac{n(\{p \in P : r_{p,s} \le t\})}{n(P)}$ or $y_s(t) = \frac{n(\{p \in P : \log_2(r_{p,s}) \le t\})}{n(P)}$,

where n(.) denotes the number of elements of a set. The performance profiles of the methods discussed in this paper are shown below.

Nwaeze et al's [32] line search method was used in all the computations since it satisfies the standard Wolfe conditions [19].

Р	FONCGM			FR			DY			PRP		
	ITE	TIM	FN.	ITE	TIM	FN.	ITE	ΤI	FN.	ITE	TIM	FN.
1	10	0.07	1.15E-	9	0.06	1.97E-12	3	0.04	2.02E-15	3	0.04	1.33E-12
2	5	0.04	6.59E-	20	0.05	1.57E-15	12	0.05	2.69E-14	20	0.04	5.16E-16
3	5	0.24	1.17E-9	20	0.35	8.86E-11	2	0.3	6.35E-10	2	0.29	4.61E-9
4	330	20.10	0.02281	336	20.56	0.02281	336	20.5	0.02281	336	20.57	0.02281
5	138	2.15	1.91666	180	2.95	1.91666	180	3.08	1.91666	175	2.55	1.91666
6	2	0.09	6.89E-	4	0.13	2.39E-10	4	0.11	2.39E-10	19	0.16	1.12E-10
7	600	3.17	2.9E-9	900	4.26	8.69E-9	900	3.34	1.78E-9	900	3.5	6.14E-10
8	761	0.86	0.29354	100	1.03	0.29354	100	0.94	0.29354	100	1.1	0.29354
9	100	46	1.93998	100	141	4.30065	100	27.3	1.67661	100	28	4.11381
1	171	0.30	1.58E-	900	0.37	2.54E-14	501	0.08	9.55E-13	900	0.34	2.33E-15
1	202	2.74	2.47E-	304	2.43	3.08E-13	323	2.81	2.39E-13	120	1.67	1.43E-12
1	9	0.35	0.05395	10	0.21	0.05395	10	0.21	0.05395	10	0.20	0.05395

Table 1. Number of iterations and CPU time in seconds

6 Discussions on Numerical Results

Table 1 contains the numerical results obtained through the new method vis-à-vis some existing methods. Table 2 displays the convergence trend of algorithm (1) on problem (12). These results indicate that the new method compares favorably well with the other methods. The execution time depends on various methods used for evaluating the step lengths and the speed of computer processing unit. We observed that the new method is relatively faster in some of the iterations recorded for the tested problems. In confirmation, Figs. 1 and 2 shows that the new method is fast and less costly as the number of function iterations per computed problem is relatively low. Finally, we saw that the results are accurate.

Table 2.	Solution	of	problem ((12)	by	algorithm	(1))
				` '	•/		•	

Iteration	f(x)	Sum of constraints
1	0.900335462627903	4.000000000000000
2	0.059832098605313	0.068911047490848
3	0.054326438636889	0.004569764836374
4	0.053979357005701	0.000052498245927
5	0.053952294603960	0.000093989215249
6	0.053950054179953	0.000036636879952
7	0.053949865266059	0.000011504805313
8	0.053949849268435	0.000004014424548
9	0.053949847897123	0.000000713781098



Fig. 1. Performance profiles on number of iterations



Fig. 2. Performance profiles on execution time

7 Conclusions

We hereby present a fourth-order nonlinear conjugate gradient method in equality constrained optimization to scientists and engineers. Some of the basic properties of the method have been explored and exploited. Numerical results show that the method is highly efficient and reliable.

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Authors' Contributions

This work was carried out in collaboration between all authors. Author NE designed the study, performed the algebraic analysis and wrote the first draft of this manuscript. Authors OG and OM did the literature searches. All authors read and approved the final manuscript.

Competing Interests

Authors have declared that no competing interests exist.

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